

Localic Relations with Open Cones

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<https://nvds.site/>

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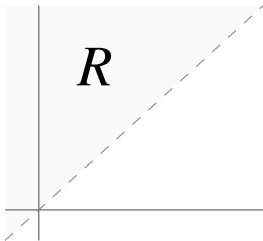


Localic relations

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spatial
relation

$$r: R \subseteq S \times S$$



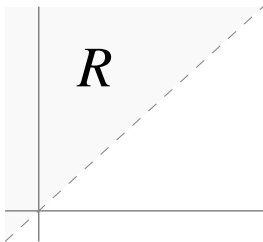
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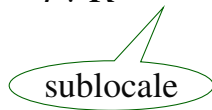
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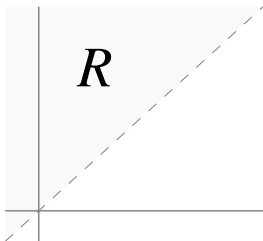
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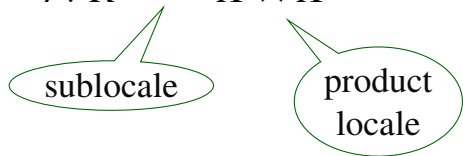
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$$\uparrow, \downarrow: \mathcal{P}(S) \longrightarrow \mathcal{P}(S)$$

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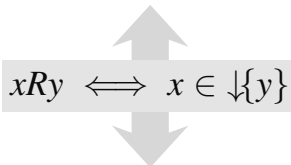
sublocale

product
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$$xRy \iff x \in \downarrow\{y\}$$

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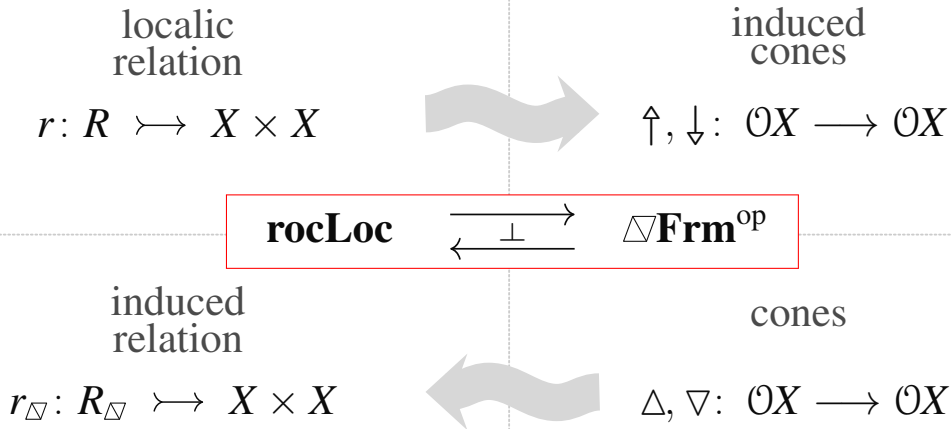
$$r_{\triangleleft}: R_{\triangleleft} \multimap X \times X$$

cones

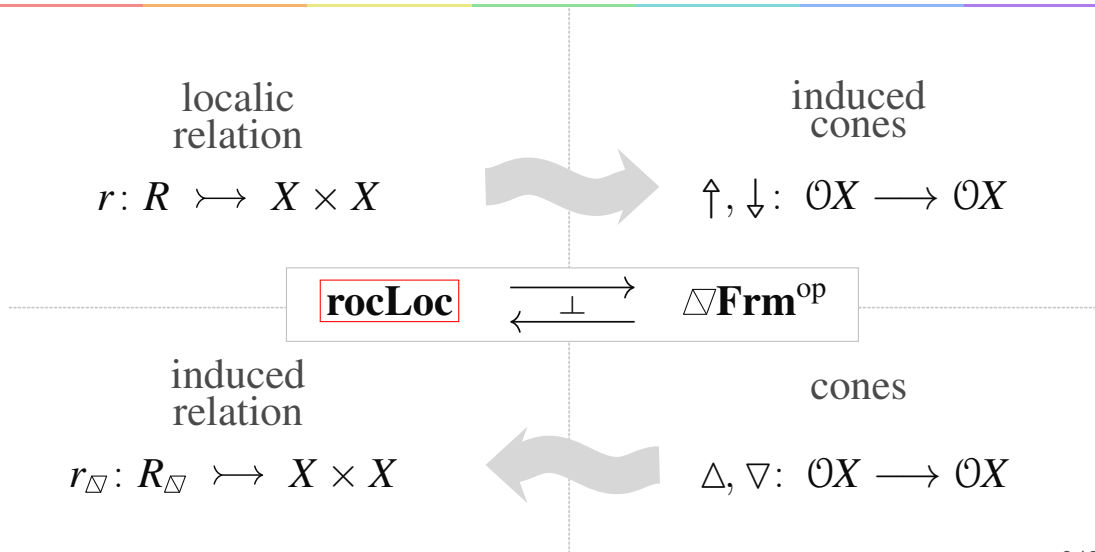
$$\Delta, \nabla: \mathcal{O}X \longrightarrow \mathcal{O}X$$



Plan



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Open cones

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$$r = (s, t): R \multimap X \times X$$

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Definition. R has *open cones*
if $s, t: R \rightarrow X$ are open.

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e.g. $(\mathbb{R}, \leq)$, interval topology, ordered vector spaces, spacetimes

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The category of open cone relations

Definition. Define a category with objects (X, R) : locales X equipped with oc R

rocLoc

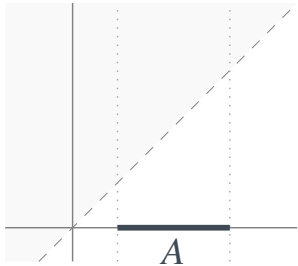
and morphisms

$$f: (X, R) \longrightarrow (Y, Q)$$

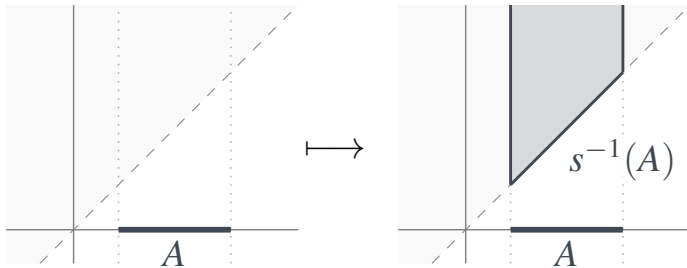
the *monotone* locale maps $f: X \rightarrow Y$:

$$(f \times f)[R] \subseteq Q$$

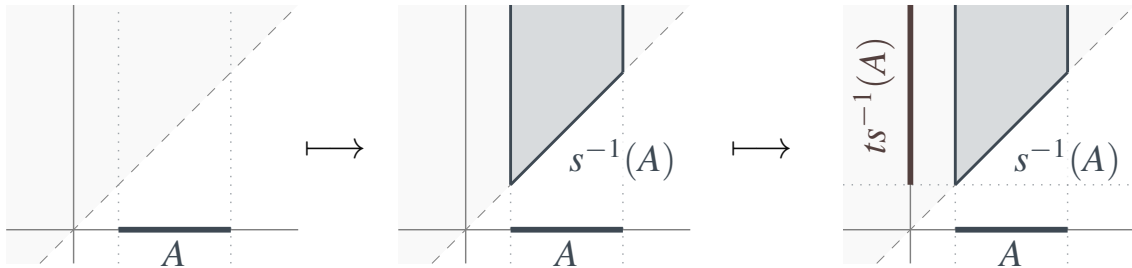
$$R \mapsto \uparrow, \downarrow$$



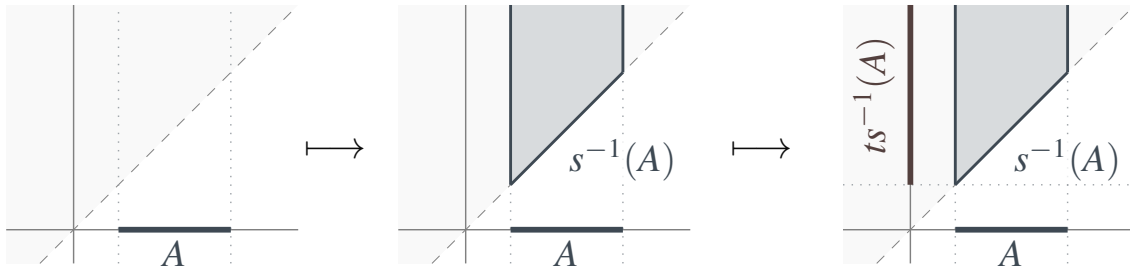
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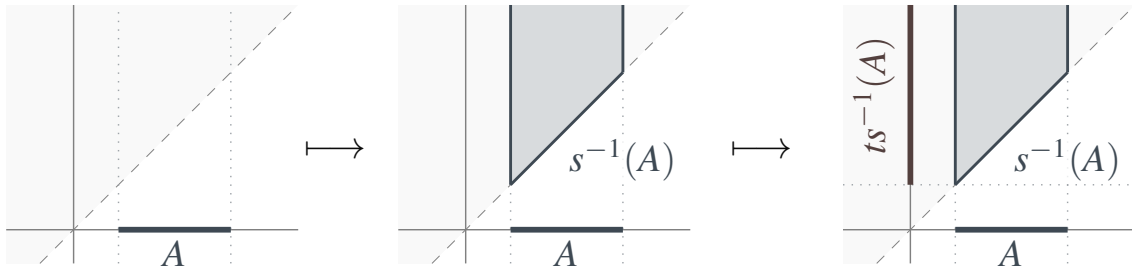


$$R \mapsto \uparrow, \downarrow$$



$$\boxed{\uparrow = ts^{-1} \quad \downarrow = st^{-1}}$$

$$R \longmapsto \uparrow, \downarrow$$



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$$\begin{aligned} \uparrow &:= t_! s^{-1} & \downarrow &:= s_! t^{-1} \\ \mathcal{O}X &\xrightarrow{s^{-1}} \mathcal{O}R \xrightarrow{t_!} \mathcal{O}X \end{aligned}$$

Conic frames?

$$\text{rocLoc} \quad \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} \quad \boxed{\triangle \mathbf{Frm}^{\text{op}}}$$

Conic frames?

$$\text{rocLoc} \quad \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad \perp \quad} \end{array} \quad \boxed{\triangle \nabla \mathbf{Frm}^{\text{op}}}$$

What should $(L, \Delta, \nabla) \in \triangle \nabla \mathbf{Frm}$ be?

Conic frames?

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What should $(L, \Delta, \nabla) \in \triangle \nabla \mathbf{Frm}$ be?

$\uparrow, \downarrow$ preserve joins, since $s!, s^{-1}, t!, t^{-1}$ do

Parallelness

When do Δ , ∇ come
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Definition. Δ, ∇ are *parallel* if:

(cf. [JT51])

$$\Delta x \wedge y \leq \Delta(x \wedge \nabla y) \quad \text{and} \quad x \wedge \nabla y \leq \nabla(\Delta x \wedge y).$$

Parallelness

When do Δ, ∇ come
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Spatially:

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Lemma. If Δ, ∇ are parallel, then:

(cf. [JT51])

$$x \wedge \nabla y = \perp \iff \Delta x \wedge y = \perp.$$

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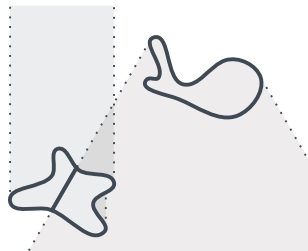
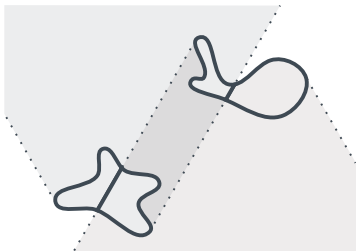
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Definition. A *conic frame* (L, Δ, ∇) is a frame L equipped with join-preserving and parallel $\Delta, \nabla: L \rightarrow L$.



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A *conic morphism*

$$h: (L, \Delta, \nabla) \longrightarrow (M, \underline{\Delta}, \bar{\nabla})$$

is a frame map $h: L \rightarrow M$:

$$\underline{\Delta} \circ h \leq h \circ \Delta,$$

$$\bar{\nabla} \circ h \leq h \circ \nabla.$$

Conic frames

$$\mathbf{cone} : \mathbf{rocLoc} \longrightarrow \triangle \mathbf{Frm}^{\mathrm{op}}$$

$$(X, R) \longmapsto (\mathcal{O}X, \uparrow, \downarrow)$$

$$f \longmapsto f^{-1}$$

Induced relation?

$$R_{\triangle} \overset{?}{\longleftarrow} (\mathcal{O}X, \Delta, \nabla)$$

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Idea: generate a congruence $\langle \sim \rangle$
on the frame $\mathcal{O}X \otimes \mathcal{O}X$

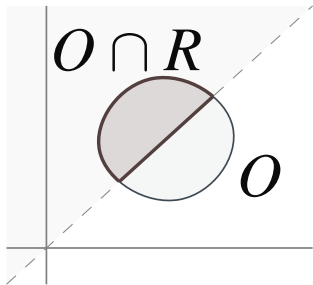
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by $\sim$ defined on rectangles $U \otimes V$
in terms of Δ, ∇

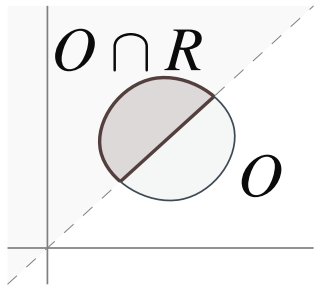
Spatial intuition



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$$r^{-1}(O) = O \cap R$$

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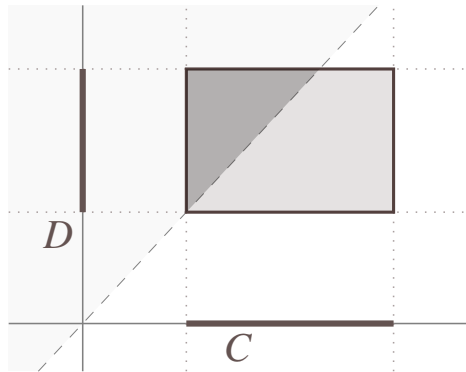
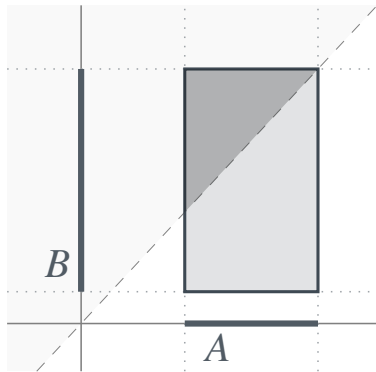


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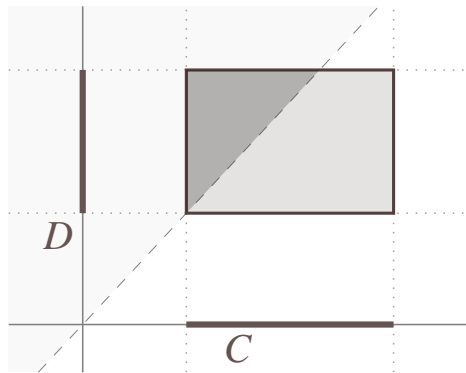
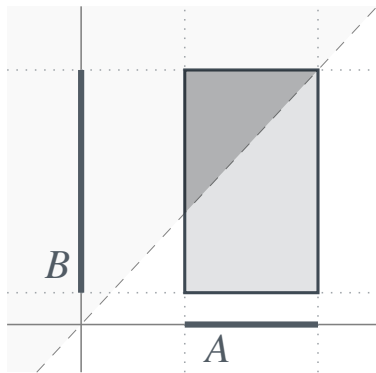
$$r^{-1}(O) = O \cap R$$

$$O \equiv O' \iff O \cap R = O' \cap R$$

Congruence on rectangles

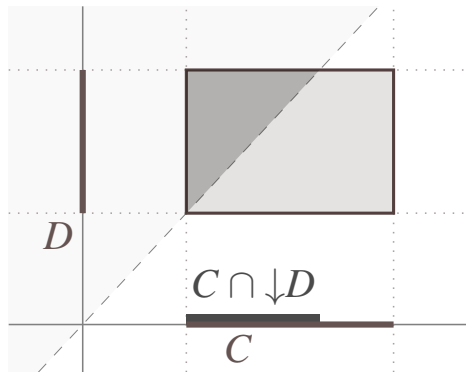
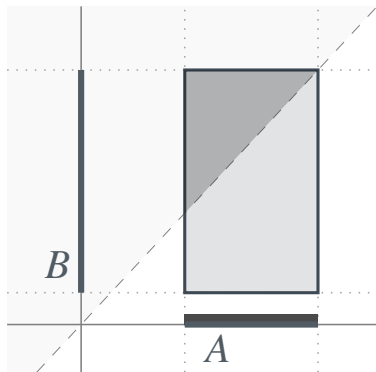


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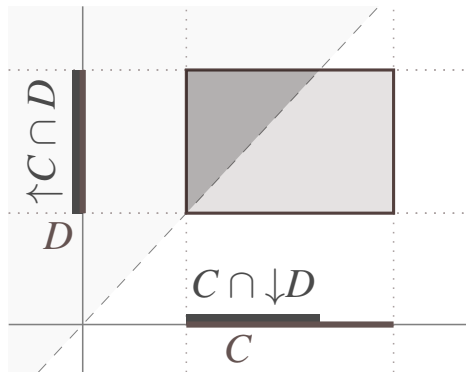
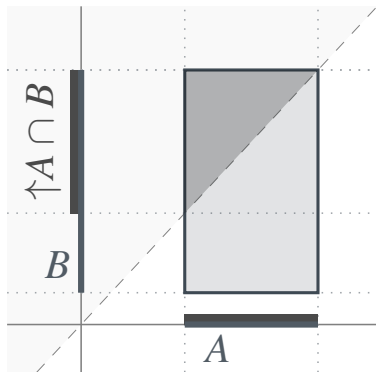
$$(A \times B) \cap R = (C \times D) \cap R \iff \text{?}$$

Congruence on rectangles



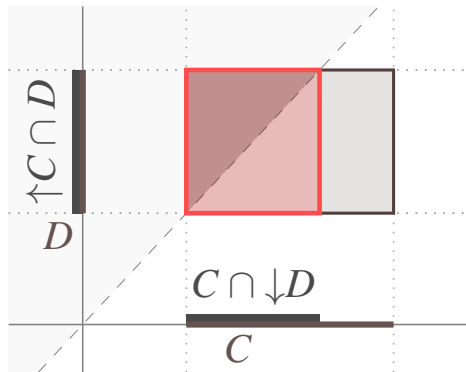
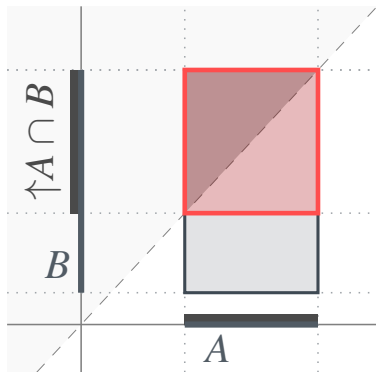
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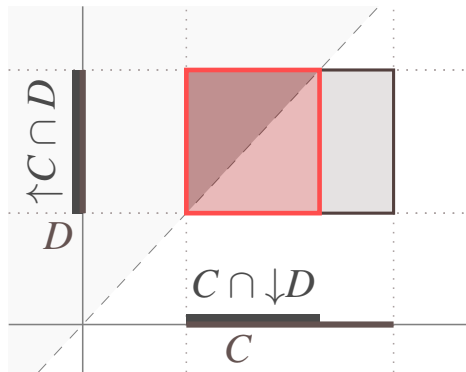
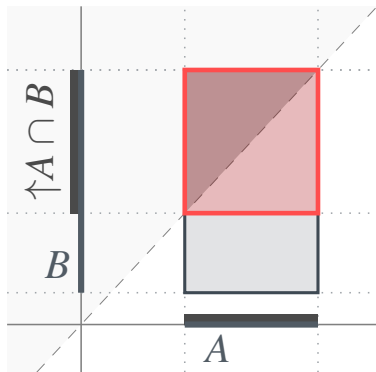
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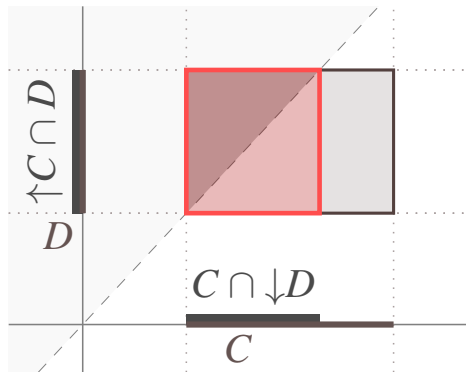
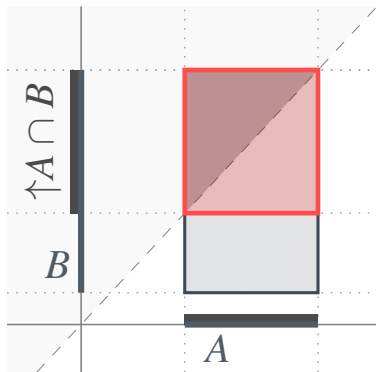
Congruence on rectangles



$$(A \times B) \cap R = (C \times D) \cap R \iff$$

$$\begin{aligned} A \cap \downarrow B &= C \cap \downarrow D \\ \uparrow A \cap B &= \uparrow C \cap D \end{aligned}$$

Congruence on rectangles



$$A \otimes B \sim C \otimes D \iff$$

$$\begin{aligned} A \wedge \nabla B &= C \wedge \nabla D \\ \Delta A \wedge B &= \Delta C \wedge D \end{aligned}$$

[MPP17] M. A. Moshier et al., “*Generating Sublocales by Subsets and Relations*”, 2017.

Definition. Given $(\mathcal{O}X, \Delta, \nabla)$,
construct R_{Δ} via:

$$\mathcal{O}R_{\Delta} := \mathcal{O}X \otimes \mathcal{O}X / \langle \sim \rangle.$$

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Definition. Given $(\mathcal{O}X, \triangle, \nabla)$,
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$$\mathcal{O}R_{\triangle} := \mathcal{O}X \otimes \mathcal{O}X / \langle \sim \rangle.$$

Theorem. $R_{\triangle}$ is an oc localic relation with

$$\uparrow = \triangle \quad \text{and} \quad \downarrow = \nabla.$$

Induced relation

$$\text{rel}: \triangleleft\!\!\triangleright\mathbf{Frm}^{\text{op}} \longrightarrow \mathbf{rocLoc}$$

$$(\mathcal{O}X, \triangle, \triangledown) \longmapsto (X, R_{\triangleleft\!\!\triangleright})$$

$$f^{-1} \longmapsto f$$

The adjunction

Theorem. There is an adjunction

$$\mathbf{rocLoc} \begin{array}{c} \xrightarrow{\text{cone}} \\ \perp \\ \xleftarrow{\text{rel}} \end{array} \triangle \mathbf{Frm}^{\text{op}}$$

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Corollary. $R \subseteq R_{\triangle}$ iff $\uparrow \leq \triangle$ and $\downarrow \leq \nabla$

Fixed points

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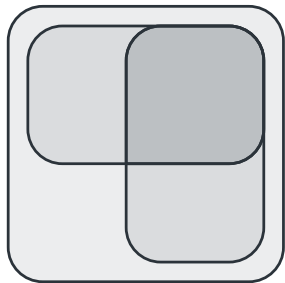
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Corollary ([Koc89]). Any weakly closed oc equivalence relation
is the kernel pair of its quotient map $X \twoheadrightarrow X/R$.

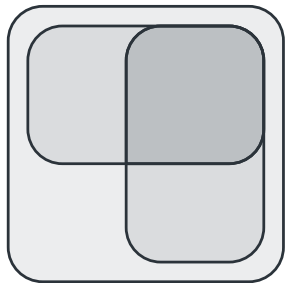
A small counterexample

$$\mathcal{OS} = \{\perp < a < \top\}$$



A small counterexample

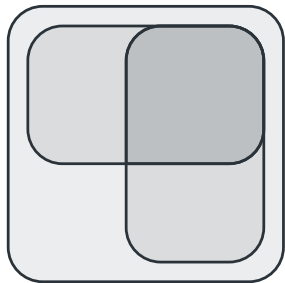
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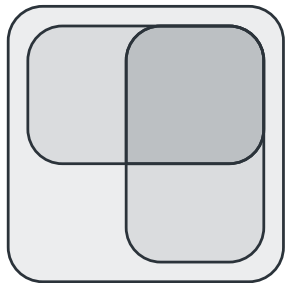


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Cones $\uparrow = \downarrow$ send
 $\perp \mapsto \perp$ and $a, \top \mapsto \top$

A small counterexample

$$\mathcal{OS} = \{\perp < a < \top\}$$



Take R open sublocale
 $(a \otimes \top) \vee (\top \otimes a)$

Cones $\uparrow = \downarrow$ send
 $\perp \mapsto \perp$ and $a, \top \mapsto \top$

$$R_{\uparrow\downarrow} = \mathbb{S} \times \mathbb{S} \supsetneq R$$

Properties in terms of cones

Proposition. If an oc R is

- reflexive then $\text{id} \leq \uparrow$ and $\text{id} \leq \downarrow$
- transitive then $\uparrow^2 \leq \uparrow$ and $\downarrow^2 \leq \downarrow$
- symmetric then $\uparrow = \downarrow$

Properties in terms of cones

Proposition. an oc $R = R_{\uparrow\downarrow}$ is

- reflexive *iff* $\text{id} \leq \uparrow$ and $\text{id} \leq \downarrow$
- transitive *iff* $\uparrow^2 \leq \uparrow$ and $\downarrow^2 \leq \downarrow$
- symmetric *iff* $\uparrow = \downarrow$

Localic Esakia duality

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$$\mathbf{CohFrm} \overset{[\text{Tow97}]}{\longleftrightarrow} \mathbf{PLoc}$$

[Tow97] C. Townsend, “*Localic Priestley Duality*”, 1997.

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Stone locales X
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$$\begin{array}{ccc} \mathbf{CohFrm} & \overset{[\text{Tow97}]}{\longleftrightarrow} & \mathbf{PLoc} \\ \cup & & \\ \mathbf{HFrm} & & \end{array}$$

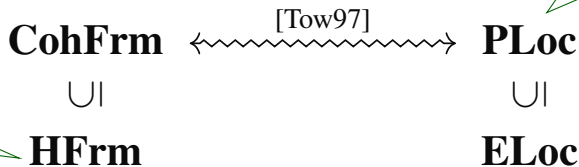
Heyting frames
[BCM23]

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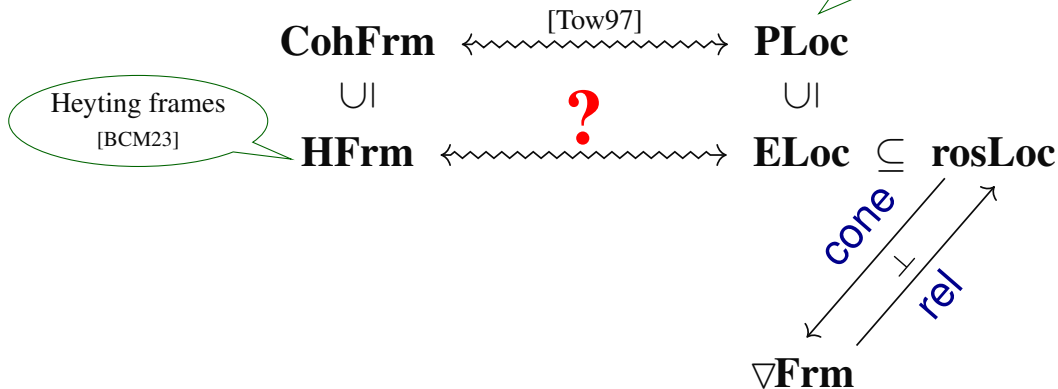
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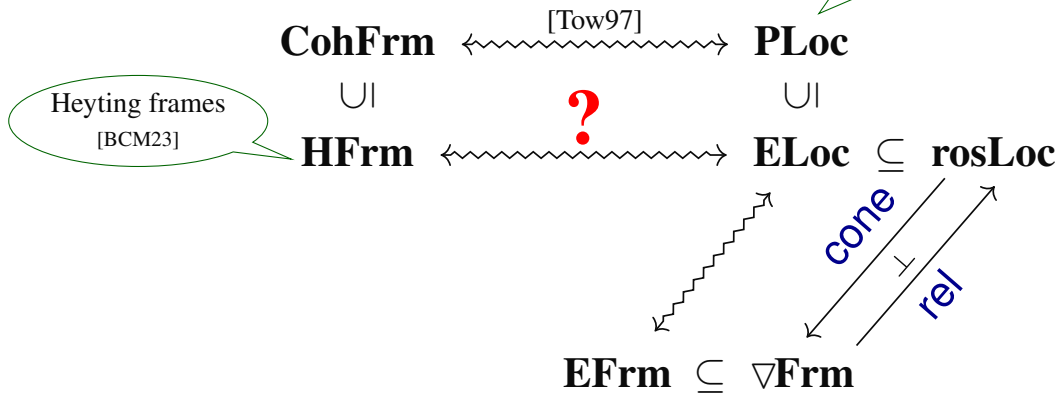


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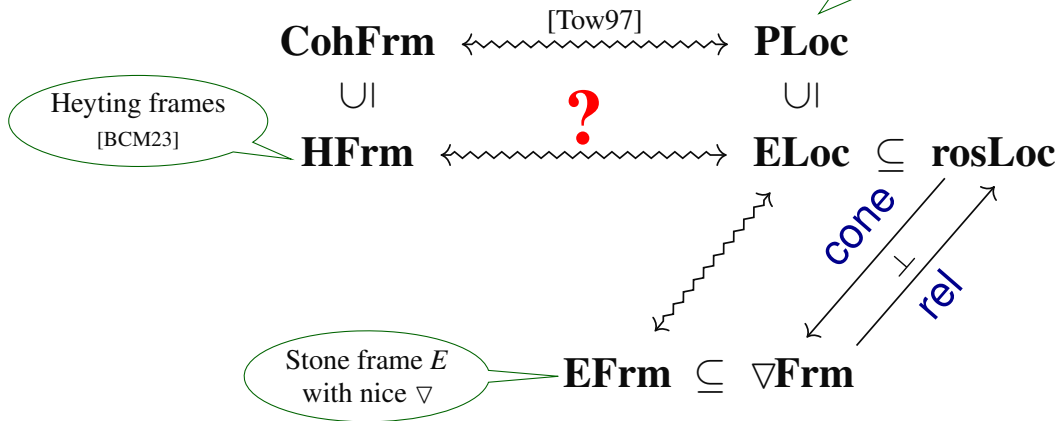


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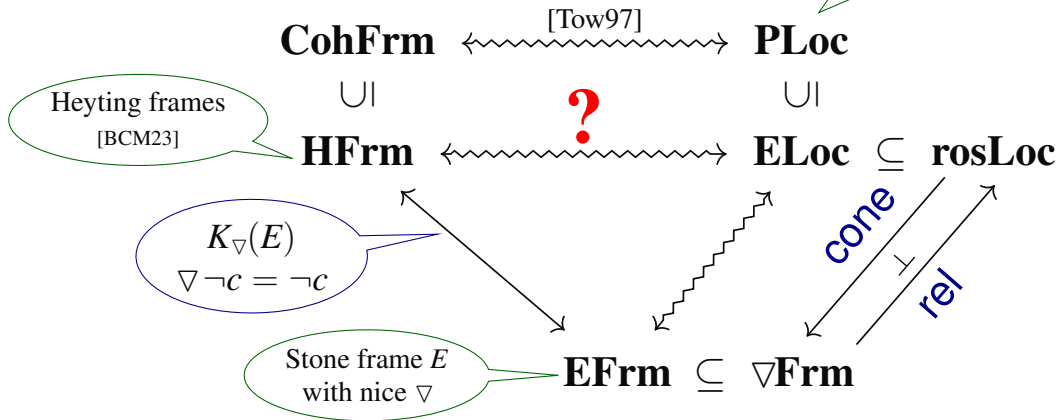


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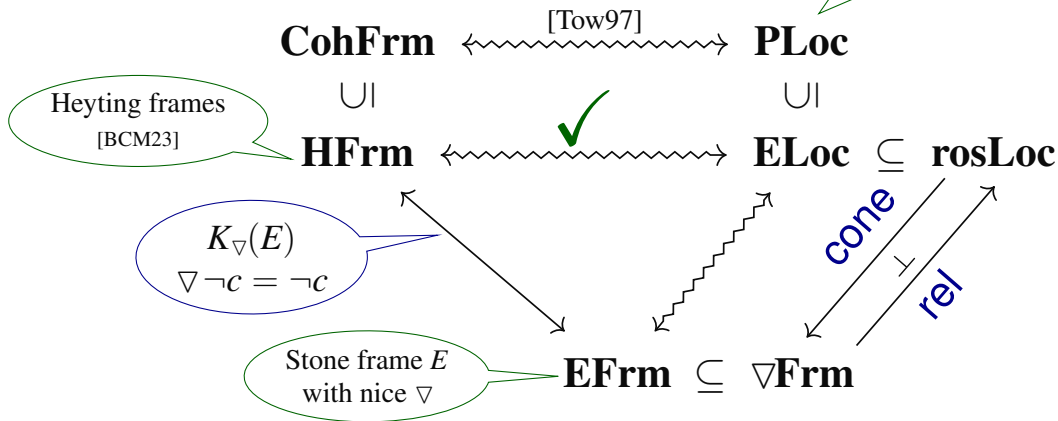


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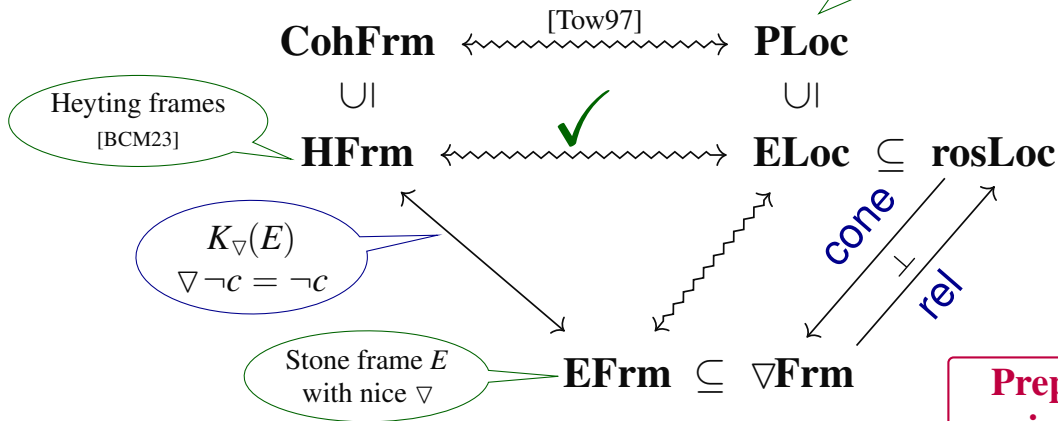


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Future work

- Recover $R \succcurlyeq X \times X$ from more general
 $\uparrow, \downarrow: \mathbf{Sl}(X) \longrightarrow \mathbf{Sl}(X)?$ (there: $\uparrow A = t[s^{-1}[A]]$ or $R \circ A$)

[Gou26] J. Goubault-Larrecq, *Stone Duality for Preordered Topological Spaces*, 2026.

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arXiv:2605.04038

Thanks!

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