

Some new semidirect-product-type constructions of residuated lattices

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- ▶ As soon as these two axioms are dropped, the known representation techniques begin to fail, or become so abstract that they say very little about the structure of the algebra itself.
- ▶ This talk reports on some progress in a project which tries to fill exactly this gap.

Our aim

- ▶ **Goal:** to introduce new product-type constructions for certain classes of residuated structures (hoops and integral residuated lattices) which behave well with respect to associativity.
- ▶ **Main representation** (finite case): $\mathbf{A} \cong \mathbf{F} \ltimes (\mathbf{A}/F)$ for every filter $F \in \mathbf{Fil} \mathbf{A}$.
- ▶ **Inspiration:** wreath and semidirect products of semigroups, and the Krohn–Rhodes theorem.
- ▶ The hoop case is where the technique is easiest to see; the point of the talk is that the same technique transfers.

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- ▶ The hoop case is where the technique is easiest to see; the point of the talk is that the same technique transfers.

If $\mathbf{A}$ and $\mathbf{B}$ are hoops and $(T_b)_{b \in |\mathbf{B}|}$ is a suitable system of transformations, we want the universe of the product to be

$$|\mathbf{A} \ltimes \mathbf{B}| = \sum_{b \in |\mathbf{B}|} T_b(\mathbf{A}),$$

i.e. every element of the “skeleton” $\mathbf{B}$ is replaced by a transformed copy of $\mathbf{A}$.

The model: wreath products

- ▶ In a wreath product the second factor plays the role of a skeleton, and every element of the skeleton induces a transformation of the first factor.
- ▶ Associativity holds because the relevant transformations are composed in the right way (as mappings): the cascade expansion does not depend on the order in which it is performed.

Hoops

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Definition

A hoop is an algebra $\mathbf{A} = (A; \cdot, \rightarrow, 1)$ of type $\langle 2, 2, 0 \rangle$ such that $(A; \cdot, 1)$ is a commutative monoid satisfying the identities

$$(H1) \quad x \rightarrow x = 1,$$

$$(H2) \quad (x \cdot y) \rightarrow z = x \rightarrow (y \rightarrow z),$$

$$(H3) \quad x \cdot (x \rightarrow y) = y \cdot (y \rightarrow x).$$

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The induced order is given by $x \leq y$ if and only if $x \rightarrow y = 1$, and moreover $x \wedge y = x \cdot (x \rightarrow y)$.

Equivalently, a hoop is a naturally ordered commutative monoid satisfying the adjointness property

$$x \cdot y \leq z \iff x \leq y \rightarrow z.$$

Filters

If $\mathbf{A} = (A; \cdot, \rightarrow, 1)$ is a hoop, then a *filter* is a nonempty subset $F \subseteq A$ such that

(F1) if $x \in F$ and $x \leq y$, then $y \in F$,

(F2) if $x, y \in F$, then $x \cdot y \in F$.

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The set of all filters of $\mathbf{A}$ is denoted by $\mathbf{Fil} \mathbf{A}$. Every filter is a subalgebra (written $\mathbf{F}$) and induces the congruence

$$\theta_F = \{\langle x, y \rangle \in A^2 \mid (x \rightarrow y) \cdot (y \rightarrow x) \in F\},$$

so that we may write $\mathbf{A}/F$ for the corresponding quotient. This correspondence between filters and congruences is a lattice isomorphism.

The f -product of hoops

Definition

Let $\mathbf{A} = (A; \cdot, \rightarrow, 1)$ and $\mathbf{B} = (B; \cdot, \rightarrow, 1)$ be hoops. A mapping $f: A \longrightarrow B$ is called a *product morphism* from $\mathbf{A}$ to $\mathbf{B}$ if it satisfies

(pM1) $f(1) = 1$,

(pM2) $f(x) \cdot f(y) = f(x \cdot y) = f(x) \wedge f(y) = f(x \wedge y)$

for all $x, y \in A$. (Hence f is monotone and $f(A) \subseteq \mathbf{Id} \mathbf{B}$.)

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Theorem

If $\mathbf{A}$ and $\mathbf{B}$ are hoops and $f: \mathbf{B} \rightarrow \mathbf{A}$ is a product morphism, then

$$\mathbf{A} \ltimes_f \mathbf{B} = (\sum_{x \in B} (f(x)], \cdot, \rightarrow, (1, 1)),$$

where

$$(\cdot) \quad (a, x) \cdot (b, y) := (a \cdot b, x \cdot y),$$

$$(\rightarrow) \quad (a, x) \rightarrow (b, y) := (f(x \rightarrow y) \wedge (a \rightarrow b), x \rightarrow y),$$

is a hoop, called the *f -product* of $\mathbf{A}$ and $\mathbf{B}$. Its order is inherited from the direct product: $(a, x) \leq (b, y)$ iff $a \leq b$ and $x \leq y$.

Two extreme examples

Example

For arbitrary hoops $\mathbf{A}, \mathbf{B}$ the constant mapping $\varepsilon: \mathbf{B} \longrightarrow \mathbf{A}$, $\varepsilon(x) = 1$, is a product morphism, and it is the greatest one with respect to the natural order. Here $\mathbf{A} \ltimes_{\varepsilon} \mathbf{B} = \mathbf{A} \times \mathbf{B}$.

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Example

If $\mathbf{A}$ has a least element 0 , then $\sigma: \mathbf{B} \longrightarrow \mathbf{A}$ given by $\sigma(1) = 1$ and $\sigma(x) = 0$ for $x \neq 1$ is the least product morphism, and $\mathbf{A} \ltimes_{\sigma} \mathbf{B} = \mathbf{B} \oplus \mathbf{A}$ (the ordinal sum).

Associated pairs of product morphisms

Definition

Let $\mathbf{A}$, $\mathbf{B}$ and $\mathbf{C}$ be hoops. A pair (f, g) is a left-associated pair of product morphisms with respect to $(\mathbf{A}, \mathbf{B}, \mathbf{C})$ if

$$f: \mathbf{B} \longrightarrow \mathbf{A} \quad \text{and} \quad g: \mathbf{C} \longrightarrow \mathbf{A} \ltimes_f \mathbf{B}$$

are product morphisms; we therefore have the hoop $(\mathbf{A} \ltimes_f \mathbf{B}) \ltimes_g \mathbf{C}$.

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Similarly, (f, g) is a right-associated pair with respect to $(\mathbf{A}, \mathbf{B}, \mathbf{C})$ if

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In a left-associated pair, g maps into the algebra $\mathbf{A} \ltimes_f \mathbf{B}$ that has already been built, so it also records how $\mathbf{C}$ acts on $\mathbf{A}$. The correspondences α and β below do nothing but rearrange this information.

Associativity up to isomorphism

Theorem

Let $\mathbf{A}, \mathbf{B}, \mathbf{C}$ be hoops.

- (i) If (f, g) is left-associated and we write $g(c) = (g_1(c), g_2(c))$, then

$$\overline{g}(c) = g_2(c), \quad \overline{f}(b, c) = f(b) \wedge g_1(c)$$

define a right-associated pair $\alpha(f, g) = (\overline{f}, \overline{g})$.

- (ii) If (f, g) is right-associated, then

$$\overline{f}(b) = f(b, 1), \quad \overline{g}(c) = (f(g(c), c), g(c))$$

define a left-associated pair $\beta(f, g) = (\overline{f}, \overline{g})$.

- (iii) α and β are mutually inverse bijections, and if $\alpha(f, g) = (\overline{f}, \overline{g})$, then

$$(\mathbf{A} \ltimes_f \mathbf{B}) \ltimes_g \mathbf{C} \cong \mathbf{A} \ltimes_{\overline{f}} (\mathbf{B} \ltimes_{\overline{g}} \mathbf{C}).$$

Decomposition of finite hoops

Lemma

Let $\mathbf{A}$ be a finite hoop and $F \in \mathbf{Fil} \mathbf{A}$. Denote by t_X the top element and by l_X the least element of the class $X \in \mathbf{A}/F$, and write l instead of l_F . Then, for all $X, Y \in \mathbf{A}/F$ and all $a \in F$:

- (i) $X \leq Y$ if and only if $t_X \leq t_Y$,
- (ii) $t_X \wedge t_Y = t_{X \wedge Y}$ and $t_X \rightarrow t_Y = t_{X \rightarrow Y}$,
- (iii) $t_X \cdot t_{X \rightarrow Y} = t_{X \wedge Y}$,
- (iv) $a \rightarrow t_X = t_X$ and $t_X \cdot a = t_X \wedge a$,
- (v) $t_X \cdot l = t_X \wedge l = l_X$.

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- (ii) $t_X \wedge t_Y = t_{X \wedge Y}$ and $t_X \rightarrow t_Y = t_{X \rightarrow Y}$,
- (iii) $t_X \cdot t_{X \rightarrow Y} = t_{X \wedge Y}$,
- (iv) $a \rightarrow t_X = t_X$ and $t_X \cdot a = t_X \wedge a$,
- (v) $t_X \cdot l = t_X \wedge l = l_X$.

Hence every class X is, via $a \mapsto a \vee l$, in bijection with the principal downset $(t_X \vee l]$ of the filter $\mathbf{F}$.

The decomposition theorem

Theorem

If $\mathbf{A}$ is a finite hoop and $F \in \mathbf{Fil} \mathbf{A}$, then $\psi: \mathbf{A}/F \longrightarrow \mathbf{F}$ defined by $\psi(X) = t_X \vee I$ is a product morphism and, moreover, $\mathbf{A} \cong \mathbf{F} \ltimes_{\psi} (\mathbf{A}/F)$.

The isomorphism is given by $\Omega(a) = (a \vee I, a/F)$, with $\Omega^{-1}(a, X) = a \wedge t_X$; the two pictures shown next are decompositions of exactly this kind.

Finite hoops as products of MV-chains

A hoop $\mathbf{A}$ is *irreducible* if $\mathbf{A} \cong \mathbf{B} \times_f \mathbf{C}$ implies that $\mathbf{B}$ or $\mathbf{C}$ is trivial.

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Theorem

A finite hoop is irreducible if and only if it is simple.

Theorem (Blok, Ferreirim)

Simple finite hoops are precisely the finite MV-chains.

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Corollary

For every finite hoop $\mathbf{A}$ there exist finite MV-chains $\mathbf{M}_{i_1}, \dots, \mathbf{M}_{i_n}$ and product morphisms $f_1, \dots, f_{n-1}$ such that

$$\mathbf{A} \cong \mathbf{M}_{i_1} \times_{f_1} (\mathbf{M}_{i_2} \times_{f_2} (\dots \times_{f_{n-2}} (\mathbf{M}_{i_{n-1}} \times_{f_{n-1}} \mathbf{M}_{i_n}) \dots)),$$

and, by associativity, also with the opposite bracketing.

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The same conclusion is known via poset products: every finite GBL-algebra, hence every finite hoop, is a poset product of finite MV-chains [Di Nola–Lettieri 2003; Jipsen–Montagna 2009]. New here is a *binary, associative* product — and that is what survives when divisibility is dropped.

Products with an MV-chain factor

Let $\mathbf{A}$ be a finite hoop and let $\mathbf{M}$ be a finite MV-chain, so that $\mathbf{Id} \mathbf{M} = \{0, 1\}$. For an idempotent $f \in \mathbf{Id} \mathbf{A}$, both

$$\psi_f: \mathbf{A} \longrightarrow \mathbf{M}, \quad \psi_f(x) = 1 \iff f \leq x,$$

$$\psi_f: \mathbf{M} \longrightarrow \mathbf{A}, \quad \psi_f(1) = 1 \quad \text{and} \quad \psi_f(x) = f \text{ for } x \neq 1,$$

are product morphisms.

Theorem

$\psi \mapsto \min \psi^{-1}(1)$ and $\psi \mapsto \psi(0)$ are bijections between $\mathbf{Id} \mathbf{A}$ and the set of all product morphisms $\mathbf{A} \longrightarrow \mathbf{M}$, resp. $\mathbf{M} \longrightarrow \mathbf{A}$.

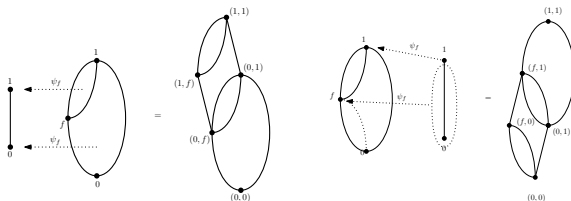


Figure 1: $\mathbf{M} \times_{\psi_f} \mathbf{A}$ and $\mathbf{A} \times_{\psi_f} \mathbf{M}$.

The f -product as a split extension

The trivial hoop is a zero object, so kernels are exactly filters. Now $\pi_2: \mathbf{A} \ltimes_f \mathbf{B} \longrightarrow \mathbf{B}$ has kernel $\pi_2^{-1}(1) = \{(a, 1) \mid a \in A\} \cong \mathbf{A}$, and it *splits*:

$$s: \mathbf{B} \longrightarrow \mathbf{A} \ltimes_f \mathbf{B}, \quad s(x) = (f(x), x), \quad \pi_2 \circ s = \text{id}_{\mathbf{B}},$$

since $s(x) \cdot s(y) = (f(x \cdot y), x \cdot y)$ and $s(x) \rightarrow s(y) = (f(x \rightarrow y), x \rightarrow y)$, using $f(x \rightarrow y) \leq f(x) \rightarrow f(y)$.

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Corollary

In a finite hoop every quotient splits canonically.

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Corollary

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So $\ltimes_f$ is a split extension in the sense of Bourn–Janelidze — with the acting maps being nuclei rather than automorphisms. The corresponding picture for exact sequences is developed in the appendix of the paper.

Generalisation I: nuclei

From now on **A**, **B**, **C** are arbitrary **integral residuated lattices**: neither commutativity nor divisibility is assumed.

A closure operator γ on **A** is a *nucleus* if $\gamma(x) \cdot \gamma(y) \leq \gamma(x \cdot y)$.

Why nuclei? The fibre $(f(x)]$ is the image of the nucleus $a \mapsto f(x) \rightarrow a$; noncommutatively, no such generator need exist.

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Lemma

*If a closure operator γ on **A** satisfies*

$$\gamma(x) \setminus \gamma(y) = \gamma(x \setminus y) \text{ (SN1)} \quad \text{or} \quad \gamma(x) / \gamma(y) = \gamma(x / y) \text{ (SN2),}$$

then it is a nucleus and $\gamma(x \wedge y) = \gamma(x) \wedge \gamma(y)$.

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then it is a nucleus and $\gamma(x \wedge y) = \gamma(x) \wedge \gamma(y)$.

A nucleus satisfying both (SN1) and (SN2) is a *strong nucleus*, and $\text{SN } \mathbf{A}$ denotes the set of all strong nuclei on **A**. Strong nuclei commute, $\alpha \circ \beta \in \text{SN } \mathbf{A}$, and $(\text{SN } \mathbf{A}; \circ, I_A)$ is a semilattice in which $\alpha \triangleleft \beta$ if and only if $\alpha A \subseteq \beta A$.

The product via nuclei

Definition (Jan Kühr)

If $\mathbf{A}$ and $\mathbf{B}$ are integral residuated lattices, then a mapping $\alpha: B \longrightarrow \text{SN } \mathbf{A}$ satisfying

$$\alpha[1] = I_A \quad \text{and} \quad \alpha[x] \circ \alpha[y] = \alpha[x \wedge y] = \alpha[x \cdot y] \quad (x, y \in B)$$

is a nuclei-morphism from $\mathbf{B}$ to $\mathbf{A}$, written $\alpha: \mathbf{B} \curvearrowright \mathbf{A}$.

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Theorem

If $\alpha: \mathbf{B} \curvearrowright \mathbf{A}$, then the algebra $\mathbf{A} \ltimes_{\alpha} \mathbf{B}$ with universe $\sum_{x \in B} \alpha[x]A$, unit $(1, 1)$ and

$$\begin{aligned}(a_1, b_1) \cdot (a_2, b_2) &= (\alpha[b_1 \cdot b_2](a_1 \cdot a_2), b_1 \cdot b_2), \\(a_1, b_1) \setminus (a_2, b_2) &= (a_1 \setminus \alpha[b_1](a_2), b_1 \setminus b_2), \\(a_2, b_2) / (a_1, b_1) &= (\alpha[b_1](a_2) / a_1, b_2 / b_1),\end{aligned}$$

is again an integral residuated lattice; its lattice operations are those of the nucleus images.

Associativity of the product via nuclei

Theorem

Let $\mathbf{A}, \mathbf{B}, \mathbf{C}$ be integral residuated lattices.

(i) For all $\alpha: \mathbf{B} \multimap \mathbf{A}$ and $\beta: \mathbf{C} \multimap \mathbf{A} \ltimes_{\alpha} \mathbf{B}$ there exist

$$\overline{\beta}: \mathbf{C} \multimap \mathbf{B} \quad \text{and} \quad \overline{\alpha}: \mathbf{B} \ltimes_{\overline{\beta}} \mathbf{C} \multimap \mathbf{A}$$

such that $(\mathbf{A} \ltimes_{\alpha} \mathbf{B}) \ltimes_{\beta} \mathbf{C} \cong \mathbf{A} \ltimes_{\overline{\alpha}} (\mathbf{B} \ltimes_{\overline{\beta}} \mathbf{C})$.

(ii) For all $\beta: \mathbf{C} \multimap \mathbf{B}$ and $\alpha: \mathbf{B} \ltimes_{\beta} \mathbf{C} \multimap \mathbf{A}$ there exist

$$\hat{\alpha}: \mathbf{B} \multimap \mathbf{A} \quad \text{and} \quad \hat{\beta}: \mathbf{C} \multimap \mathbf{A} \ltimes_{\hat{\alpha}} \mathbf{B}$$

such that $\mathbf{A} \ltimes_{\alpha} (\mathbf{B} \ltimes_{\beta} \mathbf{C}) \cong (\mathbf{A} \ltimes_{\hat{\alpha}} \mathbf{B}) \ltimes_{\hat{\beta}} \mathbf{C}$.

Moreover, $(\alpha, \beta) \mapsto (\overline{\alpha}, \overline{\beta})$ and $(\alpha, \beta) \mapsto (\hat{\alpha}, \hat{\beta})$ are mutually inverse bijections.

The transfer formulas, and the relation to f -products

The nuclei-morphisms above are given explicitly by

$$\begin{aligned}\overline{\beta}[c](b) &= (\pi_2 \circ \beta[c])(1, b), & \overline{\alpha}[b, c](a) &= (\alpha[b] \circ \beta_1[c])(a), \\ \widehat{\alpha}[b](a) &= \alpha[b, 1](a), & \widehat{\beta}[c](a, b) &= (\alpha[\beta[c](b), c](a), \beta[c](b)),\end{aligned}$$

where $\beta_1[c](a) = (\pi_1 \circ \beta[c])(a, 1)$.

The transfer formulas, and the relation to f -products

The nuclei-morphisms above are given explicitly by

$$\begin{aligned}\overline{\beta}[c](b) &= (\pi_2 \circ \beta[c])(1, b), & \overline{\alpha}[b, c](a) &= (\alpha[b] \circ \beta_1[c])(a), \\ \widehat{\alpha}[b](a) &= \alpha[b, 1](a), & \widehat{\beta}[c](a, b) &= (\alpha[\beta[c](b), c](a), \beta[c](b)),\end{aligned}$$

where $\beta_1[c](a) = (\pi_1 \circ \beta[c])(a, 1)$.

In the finite commutative case the two constructions correspond to each other one-to-one:

$$\alpha_\phi[x](a) = \phi(x) \rightarrow a, \quad \phi_\alpha(x) = \min\{a \in A \mid \alpha[x](a) = 1\}.$$

Generalisation II: product morphisms

Definition

A mapping $\phi: B \rightarrow A$ between integral residuated lattices is a product morphism if

$$\phi(1) = 1, \quad \phi(x) \wedge \phi(y) = \phi(x \wedge y), \quad \phi(x) \cdot \phi(y) \leq \phi(x \cdot y).$$

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Theorem

For a product morphism $\phi: \mathbf{B} \rightarrow \mathbf{A}$ the algebra

$$\mathbf{A} \ltimes_{\phi} \mathbf{B} = (\sum_{x \in B} (\phi(x)], \wedge, \vee, \cdot, \backslash, /, (1, 1)),$$

with componentwise $\wedge, \vee, \cdot$ and

$$(a, x) \bullet (b, y) = ((a \bullet b) \wedge \phi(x \bullet y), x \bullet y) \quad \text{for } \bullet \in \{\backslash, /\},$$

is a residuated lattice.

Associativity for integral residuated lattices

Theorem

All mappings below are product morphisms between integral residuated lattices.

(i) *For all $\alpha: \mathbf{B} \rightarrow \mathbf{A}$ and $\beta: \mathbf{C} \rightarrow \mathbf{A} \ltimes_{\alpha} \mathbf{B}$ there exist*

$$\rho\beta: \mathbf{C} \rightarrow \mathbf{B} \quad \text{and} \quad \rho\alpha: \mathbf{B} \ltimes_{\rho\beta} \mathbf{C} \rightarrow \mathbf{A}$$

such that $(\mathbf{A} \ltimes_{\alpha} \mathbf{B}) \ltimes_{\beta} \mathbf{C} \cong \mathbf{A} \ltimes_{\rho\alpha} (\mathbf{B} \ltimes_{\rho\beta} \mathbf{C})$.

(ii) *For all $\beta: \mathbf{C} \rightarrow \mathbf{B}$ and $\alpha: \mathbf{B} \ltimes_{\beta} \mathbf{C} \rightarrow \mathbf{A}$ there exist*

$$\lambda\alpha: \mathbf{B} \rightarrow \mathbf{A} \quad \text{and} \quad \lambda\beta: \mathbf{C} \rightarrow \mathbf{A} \ltimes_{\lambda\alpha} \mathbf{B}$$

such that $\mathbf{A} \ltimes_{\alpha} (\mathbf{B} \ltimes_{\beta} \mathbf{C}) \cong (\mathbf{A} \ltimes_{\lambda\alpha} \mathbf{B}) \ltimes_{\lambda\beta} \mathbf{C}$.

Here ρ and λ are mutually inverse bijections, given by the very same formulas as in the hoop case: $\rho\beta(c) = \beta_2(c)$,

$\rho\alpha(b, c) = \alpha(b) \wedge \beta_1(c)$, $\lambda\alpha(b) = \alpha(b, 1)$,

$\lambda\beta(c) = (\alpha(\beta(c), c), \beta(c))$.

Decomposition in the split case

Definition

Writing $\text{Norm}(\mathbf{A}) = \text{Cent}(\mathbf{A}) \cap \text{Id}(\mathbf{A})$, we call a residuated lattice $\mathbf{A}$ split if it is finite, distributive and integral, and if $x \cdot y = x \wedge y$ for all $x \in \text{Norm}(\mathbf{A})$ and $y \in \mathbf{A}$. The class of all such algebras is denoted by SPL .

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Lemma

Let $\mathbf{A}$ be a finite integral residuated lattice, $F \in \text{Filt}(\mathbf{A})$, and let t_X, l_X be the top and the bottom element of $X \in \mathbf{A}/F$; put $I = I_F$. Then

1. $X \leq Y$ iff $t_X \leq t_Y$ iff $l_X \leq l_Y$,
2. $t_X \wedge t_Y = t_{X \wedge Y}$, $t_X \setminus t_Y = t_{X \setminus Y}$ and $t_X / t_Y = t_{X / Y}$,
3. if $\mathbf{A} \in \text{SPL}$, then $t_X \wedge I = l_X$.

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3. if $\mathbf{A} \in \text{SPL}$, then $t_X \wedge I = l_X$.

Theorem

If $\mathbf{A} \in \text{SPL}$ and $F \in \text{Filt}(\mathbf{A})$, then $\phi(X) = t_X \vee I$ defines a product morphism $\phi: \mathbf{A}/F \rightarrow \mathbf{F}$ and $\mathbf{A} \cong \mathbf{F} \ltimes_{\phi} (\mathbf{A}/F)$.

Thank you for your attention!

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