

A logic for weighted Markov chains

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A classical Markov chain

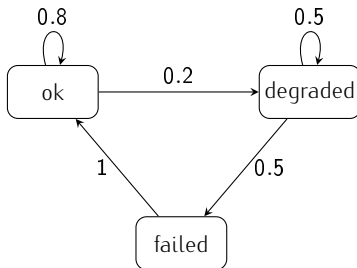
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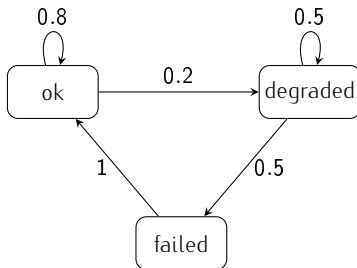
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Formally: a set of **states** X and, for each state x , a **probability distribution** $\alpha(x)$ on X giving the probability of moving to another state.

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Chains recording this data are what we'll be calling weighted Markov chains.

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The engine

A **Stone-type duality** between **weighted Markov chains** and **modal ℓ -groups**, that **extends** the Furber-Mardare-Mio duality.

Where this sits

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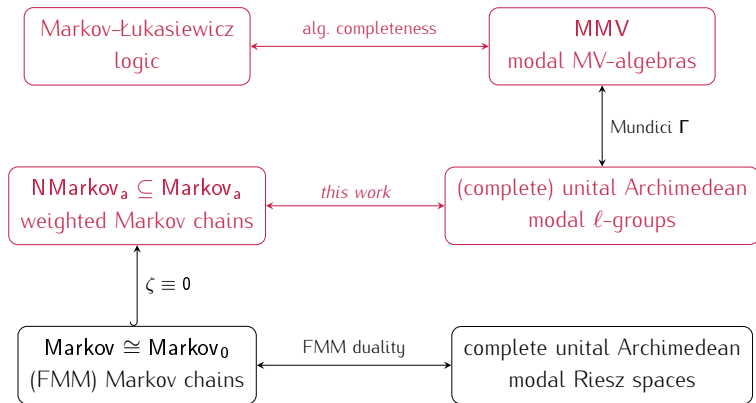
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We want a **simple** many-valued modal logic whose modality is **probabilistic** rather than normal, with a **good algebraic semantics** – and complete with respect to Markov chains.

The roadmap



Background

Weighted Markov chains

The algebraic interpretation

The logic

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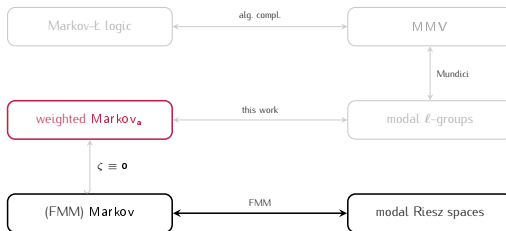
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- Mundici's Γ : unital ℓ -groups $\simeq$ MV-algebras. The functor, on objects, is $\Gamma(G, u) = [0, 1_G]$.

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An a -space (X, ζ) is (arithmetically) **normal** if it is compact and, for any two distinct points $x, y \in X$, there exist two disjoint open neighborhoods U and V of x and y respectively, such that, for every $z \in X \setminus (U \cup V)$ we have $\zeta(z) = 0$.

The main example

Let $\text{den}: \mathbb{R} \rightarrow \mathbb{N}$ be defined as $\text{den}(z) = \begin{cases} 0 & z \in \mathbb{R} \setminus \mathbb{Q} \\ \text{denominator of } z & \text{otherwise} \end{cases}$.

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Generalizing, define $\text{den}: \mathbb{R}^I \rightarrow \mathbb{N}$ as $\text{den}((p_i)_{i \in I}) = \text{lcm}\{\text{den}(p_i) \mid i \in I\}$.

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A function $f: (X, \zeta_X) \rightarrow (Y, \zeta_Y)$ is an $\mathbf{a}$ -map if it is continuous and it decreases denominators, i.e. $\zeta_Y(f(x))$ divides $\zeta_X(x)$ for all $x \in X$.

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If X is compact, $C_a(X)$ is an ℓ -group that has 1_X as strong unit, is Archimedean and is (norm-)complete.

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$C_a \dashv \text{Max}_a$

The adjunction restricts to a duality between the category of complete unital ℓ -groups and the one of normal arithmetic spaces.

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We do the same over an a-space. Let (X, ζ) be a compact Hausdorff a-space, and set

$$R_a(X) := \{F|_{C_a(X)} \mid F \in R(X)\}.$$

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- For any a-map $f: (X, \zeta_X) \rightarrow (Y, \zeta_Y)$ in the category $\mathbf{CHaus}_a$ of compact Hausdorff a-spaces, the map

$$R_a(f) : (R_a(X), \text{den}_R) \rightarrow (R_a(Y), \text{den}_R)$$

such that $R_a(f)(F)(g) := F(g \circ f)$ for all $F \in R_a(X)$ and $g \in C_a(Y)$ is an **a-map**.

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- For any a -space $(X, \zeta) \in \mathbf{CHaus}_a$ the space $R_a(X)$ is a **normal a -space**.
- For any $(X, \zeta) \in \mathbf{Norm}_a$ and for any $F \in R_a(X)$ there exists a **unique** $F' \in R(X)$ that extends F .

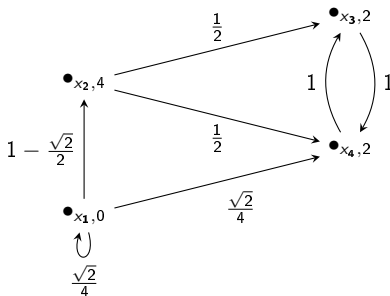
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Indeed, for any $f \in C_a(X)$ we have that $\text{den}(\alpha(x_i)(f)) | \zeta(x_i)$, since $\text{den}(f(x_i))$ is a divisor of $\zeta(x_i)$.

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So our framework **properly extends** FMM's — their category is the $\zeta \equiv 0$ case of ours.



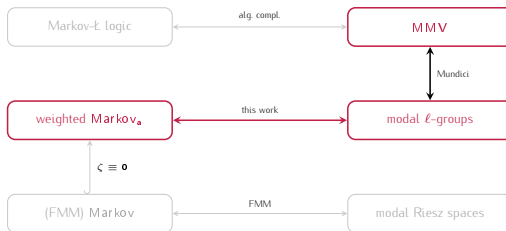
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Any weighted Markov chain induces a modal ℓ -group!

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Any weighted Markov chain induces a modal ℓ -group!

We define $C_\Diamond : \text{Markov}_a \rightarrow \text{MAulGrp}$ as follows.

- for any object $((X, \zeta), \alpha)$ in Markov_a ,
 $C_\Diamond((X, \zeta), \alpha) := (C_a(X), 1_X, \Diamond_\alpha)$, where $\Diamond_\alpha(f)(x) = \alpha(x)(f)$ for all $x \in X$,
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Morally, $\Diamond_\alpha(f)(x)$ is the expected value of f after one step from x .

From groups to chains (the idea)

Given an **Archimedean** modal unital ℓ -group $(G, 1_G, \Diamond)$ and $x \in \text{Max}_a(G) \subseteq \mathbb{R}^G$, the map $x \circ \Diamond$ is a positive homomorphism with $(x \circ \Diamond)(1_G) \leq 1$.

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Via the Riesz representation and Stone–Weierstrass theorems, it yields a **transition** $\alpha_\Diamond(x) \in R_a(\text{Max}_a(G))$, hence an **weighted Markov chain** $(\text{Max}_a(G), \alpha_\Diamond)$.

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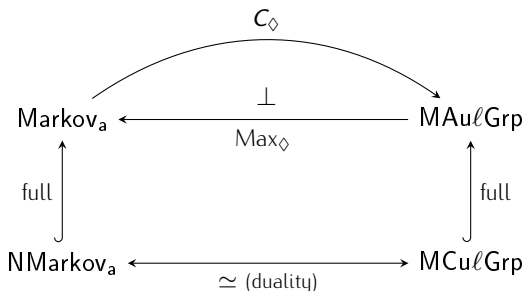
Via the Riesz representation and Stone–Weierstrass theorems, it yields a **transition** $\alpha_\diamond(x) \in R_a(\text{Max}_a(G))$, hence an **weighted Markov chain** $(\text{Max}_a(G), \alpha_\diamond)$.

This defines a functor $\text{Max}_\diamond : \mathbf{MAu\ell Grp} \rightarrow \mathbf{Markov}_a$ (pre-composition on arrows).

[more detail](#)

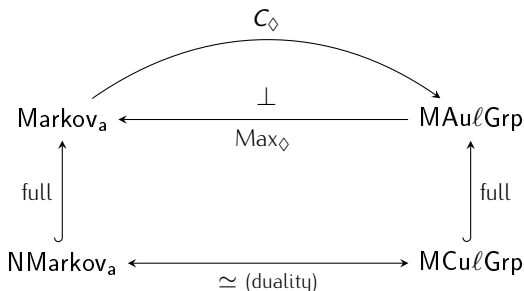
The duality

The functors $C_\diamond$ and $\text{Max}_\diamond$ form a contravariant adjunction $C_\diamond \dashv \text{Max}_\diamond$, which **restricts to a duality**:



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Proof idea: the unit and counit of the Abbadini–Marra–Spada adjunction are arrows in our categories.

Modal MV-algebras: the bridge to a logic

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A **modal MV-algebra** (or *MMV-algebra*) is an algebra $(\mathbf{A}, \Diamond)$, where $\mathbf{A}$ is an MV-algebra and $\Diamond: \mathbf{A} \rightarrow \mathbf{A}$ is a unary operation such that:

1. $\Diamond(0) = 0$,
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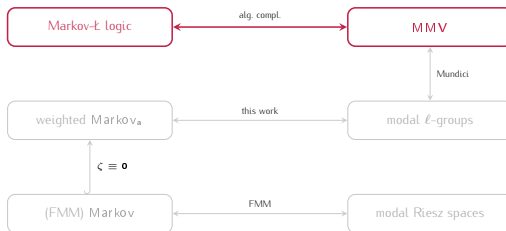
So our weighted chains carry a propositional, many-valued, modal **syntax**.

Background

Weighted Markov chains

The algebraic interpretation

The logic



Markov-Łukasiewicz logic

In the language $\mathcal{L}$ given by **Var**, the connectives of Łukasiewicz logic, and a unary connective $\Diamond$, the set **MFm** of well-formed formulas is defined recursively as follows:

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- Rules:

(MP) from φ and $\varphi \rightarrow \psi$ derive ψ ,

(M) from $(\varphi \rightarrow \psi)$ derive $(\Diamond(\varphi) \rightarrow \Diamond(\psi))$.

Algebraic completeness

Let $\Gamma \cup \{\varphi\} \subseteq \text{MFm}$. Then,

Theorem 1 (algebraic completeness)

$$\Gamma \vdash_{\text{ML}} \varphi \text{ iff } \Gamma \models_{\text{MMV}} \varphi$$

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Moreover, adding (M) as an axiom does not work!

Indeed

$$\not\models_{\mathbf{MMV}} (\varphi \rightarrow \psi) \rightarrow (\Diamond\varphi \rightarrow \Diamond\psi).$$

An example

Take $\varphi := \top$ and $\psi := p \in \text{Var}$.

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Then $v(\diamond(p)) = \diamond(v(p)) = \diamond\left(\frac{1}{2}, \frac{2}{5}\right) = \left(\frac{12}{25}, \frac{2}{5}\right)$ and

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On the other hand:

$$v(\diamond(\top) \rightarrow \diamond(p)) = (1, 1) \rightarrow \left(\frac{12}{25}, \frac{2}{5} \right) = \left(\frac{12}{25}, \frac{2}{5} \right).$$

Therefore,

$$v((\top \rightarrow p) \rightarrow (\diamond(\top) \rightarrow \diamond(p))) = \left(\frac{49}{50}, 1 \right) \neq (1, 1).$$

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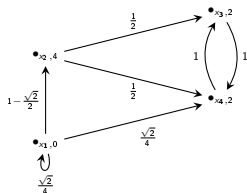
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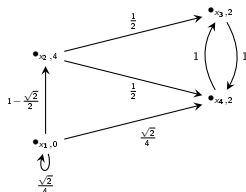
Valuation in a Markov model: an example

Take the weighted Markov chain $((X, \zeta), \alpha)$ from before, and $p \in \mathbf{Var}$ with $f_p(x_1) = 0$, $f_p(x_2) = 1$, $f_p(x_3) = 0$, $f_p(x_4) = 1$.



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$t_{\Diamond p}^{\mathfrak{M}}(x) = \alpha(x)(f_p)$ is the **expected value of p** after one step:

$$t_{\Diamond p}(x_1) = \frac{\sqrt{2}}{4} \cdot 0 + \left(1 - \frac{\sqrt{2}}{2}\right) \cdot 1 + \frac{\sqrt{2}}{4} \cdot 1 = 1 - \frac{\sqrt{2}}{4},$$

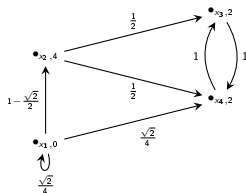
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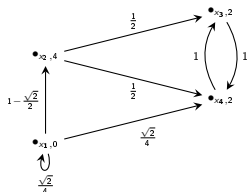
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It is an **a-map**: irrational at x_1 ($\text{den} = 0 = \zeta(x_1)$), and $\text{den}(\frac{1}{2}) = 2 \mid 4 = \zeta(x_2)$.

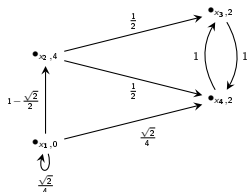
Valuation in a Markov model: an example (cont.)

Recall $f_p = (0, 1, 0, 1)$ and $t_{\Diamond p} = (1 - \frac{\sqrt{2}}{4}, \frac{1}{2}, 1, 0)$ on (x_1, x_2, x_3, x_4) .



Valuation in a Markov model: an example (cont.)

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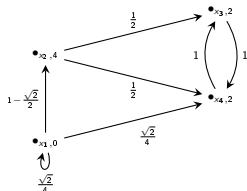


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So $\mathfrak{M}, x_2 \models \Diamond p \rightarrow p$, but $\mathfrak{M} \not\models \Diamond p \rightarrow p$ (it fails e.g. at x_3).

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So, consider the extension **aML** of the Markov-Łukasiewicz logic obtained by adding the following rule:

(A) from $n\varphi \rightarrow \psi$ for all $n \in \mathbb{N}$ derive $\neg\varphi$.

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The main point of this proof is that the Lindenbaum-Tarski algebra of $\mathbf{aML}$ is Archimedean.

This allows us to use the adjunction between $\mathbf{aMMV}$ and $\mathbf{Markov}_a$ to prove the right-to-left direction.

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In this case, we show that there exists an **adjunction** between NMarkov_a and Markov and get the result from the Corollary above.

Summing up

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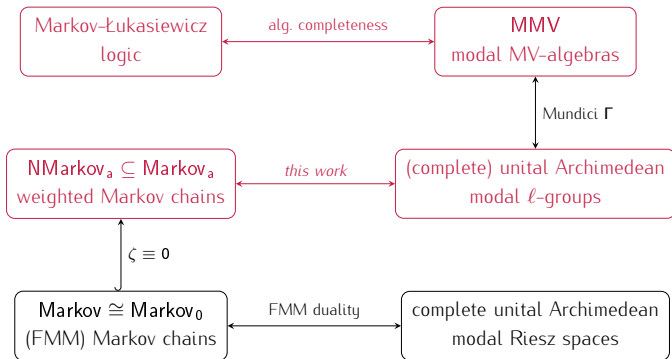
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Put it differently: is the theory of modal Riesz spaces a conservative extension of the theory of modal ℓ -groups?

Thank you for your attention!



Backup: from groups to chains (full construction)

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It is not difficult (but not trivial) to see that $\alpha_\Diamond$ is **continuous** and **decreases denominators**; hence $(\text{Max}_a(G), \alpha_\Diamond)$ is a weighted Markov chain. [back](#)