

Derivations. Briefly: use anything from propositional logic + new rules (that we discuss one by one).
The first such new rule is:

$\forall Out$: If x is **free** in $\varphi(x)$, then

$$\frac{\forall x \varphi(x)}{\varphi[x/t]}$$

$\exists In$: If $\varphi[x/t]$ is a substitution instance of $\varphi(x)$, then

$$\frac{\varphi[x/t]}{\exists x \varphi(x)}$$

(1) Apply the $\exists In$ rule in all possible ways to the formulas below (if possible).

- (a) $F(b)$
- (b) $R(c, d)$
- (c) $F(c) \wedge G(c)$

(2) Derive

$$\begin{array}{c} \frac{\forall x(F(x) \rightarrow H(x)) \quad F(a)}{\exists xH(x)} \\ \hline \frac{\forall x(G(x) \rightarrow H(x)) \quad G(b)}{\exists x(G(x) \wedge H(x))} \\ \hline \frac{\exists x \sim R(x, a) \rightarrow \sim \exists xR(a, x) \quad \sim R(a, a)}{\sim R(a, b)} \end{array}$$

K :

$$\frac{\forall x(\varphi(x) \rightarrow \psi(x))}{\forall x\varphi(x) \rightarrow \forall x\psi(x)}$$

Gen : If x is *not* free in φ , then

$$\frac{\varphi}{\forall x\varphi(x)}$$

(3) Derive

$$\begin{array}{c} \frac{\forall x(F(x) \rightarrow G(x)) \quad \forall xF(x)}{\forall xG(x)} \\ \hline \frac{\forall x(F(x) \rightarrow G(x)) \quad \forall x(G(x) \rightarrow H(x))}{\forall x(F(x) \rightarrow H(x))} \end{array}$$

(4) Which of the following formulas are logical truth?

- (a) $\exists x \forall y \varphi \rightarrow \forall y \exists x \varphi$
- (b) $\forall x \exists y \varphi \rightarrow \exists y \forall x \varphi$
- (c) $\forall x \exists x \varphi \rightarrow \exists x \forall x \varphi$